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Quantum Reed-Solomon Codes

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Overview

- Quantum information
- Quantum operations
- Quantum error-correcting codes (QECC)
- Construction of QECC
- Quantum Reed-Solomon codes
- Quantum implementation of linear transforms
- DFT over finite fields
- Concluding remarks

Quantum Information

Quantum-bit (qubit)

basis states:

$$\text{"0"} \hat{=} |0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \in \mathbb{C}^2, \quad \text{"1"} \hat{=} |1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \in \mathbb{C}^2$$

general state:

$$|q\rangle = \alpha|0\rangle + \beta|1\rangle \quad \text{where } \alpha, \beta \in \mathbb{C}, |\alpha|^2 + |\beta|^2 = 1$$

Quantum register

basis states:

$$|b_1\rangle \otimes \dots \otimes |b_n\rangle =: |b_1 \dots b_n\rangle = |\mathbf{b}\rangle \quad \text{where } b_i \in \{0, 1\}$$

general state:

$$|\psi\rangle = \sum_{\mathbf{x} \in \{0,1\}^n} c_{\mathbf{x}} |\mathbf{x}\rangle \quad \text{where } \sum_{\mathbf{x} \in \{0,1\}^n} |c_{\mathbf{x}}|^2 = 1$$

\implies normalised vector in $(\mathbb{C}^2)^{\otimes n} \cong \mathbb{C}^{2^n}$

\implies note the similarity to the group ring $\mathbb{C}[\mathbb{F}_2^n]$

Quantum Operations: "Gates"

NOT gate

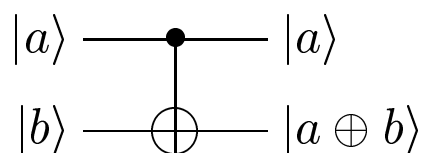
$$NOT = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \quad \begin{array}{l} |0\rangle \mapsto |1\rangle \\ |1\rangle \mapsto |0\rangle \end{array}$$

Hadamard transformation

$$H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \quad \begin{array}{l} |0\rangle \mapsto \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \\ |1\rangle \mapsto \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) \end{array}$$

Controlled *NOT* (*CNOT*) gate

$$CNOT = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \quad \begin{array}{l} |00\rangle \mapsto |00\rangle \\ |01\rangle \mapsto |01\rangle \\ |10\rangle \mapsto |11\rangle \\ |11\rangle \mapsto |10\rangle \end{array}$$



Simple Quantum Code: Bit-Flips

$$\left. \begin{array}{l} |0\rangle \mapsto |000\rangle \\ |1\rangle \mapsto |111\rangle \end{array} \right\} \implies \alpha|0\rangle + \beta|1\rangle \mapsto \alpha|000\rangle + \beta|111\rangle$$

two-dimensional subspace $\mathcal{H}_C \leq \mathcal{H}_2 \otimes \mathcal{H}_2 \otimes \mathcal{H}_2$

bit-flip	quantum state	subspace
no error	$\alpha 000\rangle + \beta 111\rangle$	$(\mathbb{1} \otimes \mathbb{1} \otimes \mathbb{1})\mathcal{H}_C$
1st position	$\alpha 100\rangle + \beta 011\rangle$	$(\sigma_x \otimes \mathbb{1} \otimes \mathbb{1})\mathcal{H}_C$
2nd position	$\alpha 010\rangle + \beta 101\rangle$	$(\mathbb{1} \otimes \sigma_x \otimes \mathbb{1})\mathcal{H}_C$
3rd position	$\alpha 001\rangle + \beta 110\rangle$	$(\mathbb{1} \otimes \mathbb{1} \otimes \sigma_x)\mathcal{H}_C$

orthogonal decomposition: $\mathcal{H}_2 \otimes \mathcal{H}_2 \otimes \mathcal{H}_2 =$

$$\begin{aligned} & (\mathbb{1} \otimes \mathbb{1} \otimes \mathbb{1})\mathcal{H}_C \\ & \oplus (\sigma_x \otimes \mathbb{1} \otimes \mathbb{1})\mathcal{H}_C \\ & \oplus (\mathbb{1} \otimes \sigma_x \otimes \mathbb{1})\mathcal{H}_C \\ & \oplus (\mathbb{1} \otimes \mathbb{1} \otimes \sigma_x)\mathcal{H}_C \end{aligned}$$

Quantum Errors

Bit-flip:

- interchange of the basis states $|0\rangle$ and $|1\rangle$
- corresponds to “classical” bit error

- *NOT* gate resp. $\sigma_x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$

Phase-flip:

- inversion of the *relative* phase of the states $|0\rangle$ and $|1\rangle$
- no classical analogue

- $\sigma_z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$

Combination:

- $\sigma_y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix} = i\sigma_x\sigma_z$

Conjugation:

- $H\sigma_xH = \sigma_z$
 - $H\sigma_zH = \sigma_x$
- } interchange of bit-flip and phase-flip

Quantum Codes (QECC)

(CALDERBANK & SHOR/STEANE)

Definition:

Let $C = [n, k, d]$ be a weakly self-dual linear binary code, i. e., $C \leq C^\perp$, where $C^\perp = [n, n - k, d^\perp]$.

For coset representatives $w \in C^\perp/C$, define

$$|\psi_w\rangle := \frac{1}{\sqrt{|C|}} \sum_{c \in C} |c + w\rangle$$

or equivalently (Hadamard transformation)

$$|\hat{\psi}_w\rangle := \frac{1}{\sqrt{|C^\perp|}} \sum_{c \in C^\perp} (-1)^{c \cdot w} |c\rangle$$

\implies quantum error-correcting code $C = [[n, n - 2k, d' \geq d^\perp]]$

Encoded quantum state:

$$|\psi\rangle = \sum_{w \in C^\perp/C} \alpha_w |\psi_w\rangle = \sum_{c \in C^\perp} \beta_c |c\rangle$$

$$|\hat{\psi}\rangle = \sum_{w \in C^\perp/C} \hat{\alpha}_w |\hat{\psi}_w\rangle = \sum_{c \in C^\perp} \gamma_c |c\rangle$$

\implies superpositions of codewords of C^\perp

Weakly Self-Dual Codes (I)

Construction of QECC:

Search for good linear binary weakly self-dual codes.

Dual basis:

For a basis $\mathcal{B} = (b_1, \dots, b_\ell)$ of $\mathbb{F}_{2^\ell} | \mathbb{F}_2$ let $\mathcal{B}^\perp = (b'_1, \dots, b'_\ell)$ denote the *dual basis* with

$$\text{tr}(b_i \cdot b'_j) = \delta_{ij}.$$

Self-dual basis:

$$\mathcal{B} = \mathcal{B}^\perp, \text{ i. e., } \text{tr}(b_i \cdot b_j) = \delta_{ij}$$

Subfield expansion:

For $C = [n, k, d] \leq \mathbb{F}_{2^\ell}^n$, define

$$\mathcal{B}(C) := \left\{ (c_{11}, \dots, c_{1\ell}, \dots, c_{n1}, \dots, c_{n\ell}) : \right. \\ \left. \mathbf{c} = (c_1, \dots, c_n) \in C, c_i = \sum_{j=1}^{\ell} c_{ij} b_j \right\}$$

$$\implies \mathcal{B}(C) = [\ell n, \ell k, d' \geq d] \leq \mathbb{F}_2^{\ell n}$$

Weakly Self-Dual Codes (II)

The following diagram commutes:

$$\begin{array}{ccc} C & \longrightarrow & C^\perp \\ \text{basis } \mathcal{B} \downarrow & & \downarrow \text{dual basis } \mathcal{B}^\perp \\ \mathcal{B}(C) & \longrightarrow & \mathcal{B}^\perp(C^\perp) = (\mathcal{B}(C))^\perp \end{array}$$

$C = [n, k] \leq \mathbb{F}_{2^\ell}$ weakly self-dual, i. e., $C \leq C^\perp$,

\mathcal{B} a self-dual basis of $\mathbb{F}_{2^\ell} | \mathbb{F}_2$

$\implies \mathcal{B}(C) \leq \mathcal{B}(C^\perp)$.

Idea:

- start with a weakly self-dual code C over \mathbb{F}_{2^ℓ} and a self-dual basis \mathcal{B}
- obtain a weakly self-dual binary code $\mathcal{B}(C)$
- obtain a CSS code

Advantages:

- correction of burst errors
- suitable for concatenation
- applies to Reed-Solomon codes



Quantum Reed-Solomon Codes

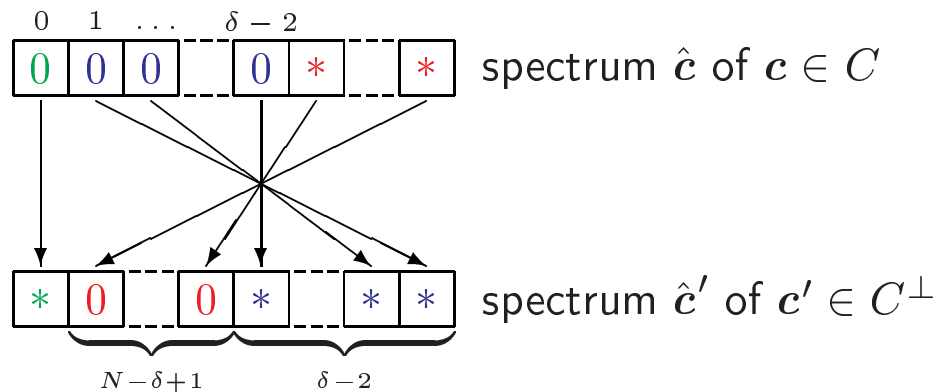
RS code $C = [n = 2^\ell - 1, n - \delta + 1, \delta]$ over \mathbb{F}_{2^ℓ} ,
generator polynomial

$$g(X) = (X - 1)(X - \alpha) \dots (X - \alpha^{\delta-2})$$

weakly self-dual iff $\delta > n/2 + 1$:

$$\begin{aligned} g^\perp(X) &= (X - \alpha^{-(\delta-1)})(X - \alpha^{-\delta}) \dots (X - \alpha^{-(n-1)}) \\ &= (X - \alpha^1)(X - \alpha^2) \dots (X - \alpha^{n-\delta+1}) \end{aligned}$$

Spectra:



Encoding & decoding:

discrete Fourier transformation over \mathbb{F}_{2^ℓ}

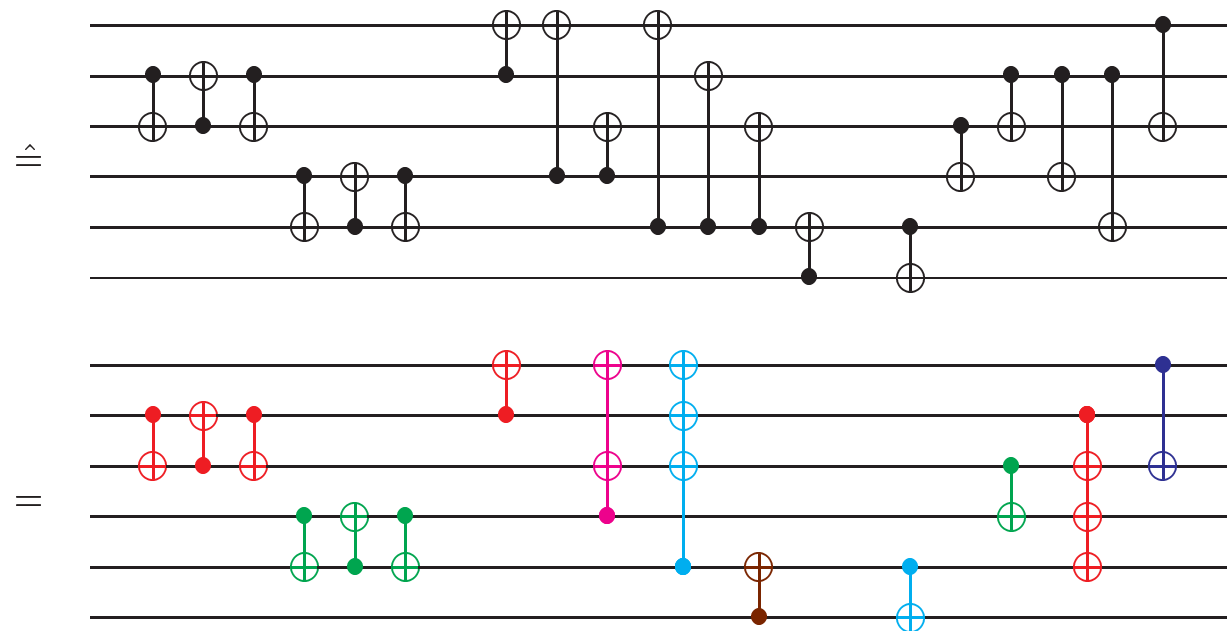
\implies linear transformation over \mathbb{F}_{2^ℓ}

choose a basis \mathcal{B} of $\mathbb{F}_{2^\ell} | \mathbb{F}_2$

\implies linear transformation over \mathbb{F}_2

Example: Linear Transformation over \mathbb{F}_2

$$\begin{pmatrix} 1 & \cdot & 1 & \cdot & \cdot & \cdot \\ \cdot & \cdot & 1 & 1 & \cdot & \cdot \\ 1 & 1 & \cdot & 1 & 1 & \cdot \\ 1 & 1 & 1 & \cdot & \cdot & 1 \\ 1 & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & 1 & \cdot \end{pmatrix} = \begin{pmatrix} 1 & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & 1 & \cdot & \cdot & \cdot \\ \cdot & 1 & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & 1 & \cdot \\ \cdot & \cdot & \cdot & 1 & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & 1 \end{pmatrix} \begin{pmatrix} 1 & \cdot & \cdot & \cdot & \cdot & \cdot \\ 1 & 1 & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & 1 & \cdot & \cdot & \cdot \\ 1 & \cdot & 1 & 1 & \cdot & \cdot \\ 1 & 1 & 1 & \cdot & 1 & \cdot \\ \cdot & \cdot & \cdot & \cdot & 1 & 1 \end{pmatrix} \begin{pmatrix} 1 & \cdot & 1 & \cdot & \cdot & \cdot \\ \cdot & 1 & 1 & 1 & 1 & \cdot \\ \cdot & \cdot & 1 & 1 & \cdot & \cdot \\ \cdot & \cdot & \cdot & 1 & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & 1 & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & 1 \end{pmatrix}$$



Conclusion & Outlook

- new method to construct QECC
- also applicable for BCH codes (see [quant-ph/9910060](#))
- correction of burst errors
- decoding either as codes over \mathbb{F}_{2^ℓ} or \mathbb{F}_2
- a lot of encoding/decoding techniques exist, e. g., using quantum linear shift registers (see [quant-ph/9910061](#))

Further references:

<http://iaks-www.ira.uka.de/home/grassl/>

<http://iaks-www.ira.uka.de/QIV/>

and a lot of preprints at

<http://xxx.lanl.gov/archive/quant-ph/>